Lotto: Secure Participant Selection against Adversarial Servers in Federated Learning

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Growth of edge computing

Edge devices generate massive **data**





Growth of edge computing



Growth of edge computing



Privacy-Enhancing
TechniqueFederated Learning!Privacy GuaranteeData kept on premises

¹McMahan et al.''Communication-Efficient Learning of Deep Networks from Decentralized Data'', In AISTATS '17

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Privacy-Enhancing Federated Learning¹ Technique **Privacy Guarantee** Data kept on premises

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2. Local training \rightarrow Local model update

Privacy-Enhancing Federated Learning¹ Technique **Privacy Guarantee** Data kept on premises

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3. Model aggregation \rightarrow Global model update

Privacy-Enhancing Federated Learning¹ Technique Data kept on premises **Privacy Guarantee**

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3. Model aggregation \rightarrow Global model update



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Ground truth



Reconstructed

Problem: Data can be reconstructed from **local model updates**²

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Privacy-Enhancing Technique

Privacy Guarantee

Data kept on premises

Federated Learning¹

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Secure Aggregation^{3,4}

Local updates unseen







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Secure Aggregation^{3,4}

Local updates unseen

Problem: Data still has footprints in global model update⁵

⁻S '17 ⁵Nasr et al.''Comprehensive Privacy Analysis of Deep Learning: Passive and Active White-box Inference Attacks against Centralized and Federated Learning'', In S&P '19



Privacy-Enhancing Technique

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⁵S '17 ⁵Nasr et al.''Comprehensive Privacy Analysis of Deep Learning: Passive and Active White-box Inference Attacks against Centralized and Federated Learning'', In S&P '19 ⁶Cynthia.''Differential Privacy'', 06.





¹Kairouz et al. "The Distributed Discrete Gaussian Mechanism for Federated Learning with Secure Aggregation", In ICML '21

²Agarwal. ''The Skellam Mechanism for Differentially Private Federated Learning'', In NeurIPS '21

Privacy-Enhancing Technique	Federated Learning ¹	Secure Aggregation	Differential Privacy
Privacy Guarantee	Data kept on premises	Local updates unseen	Global update leaks little about any client

May **not** hold

Privacy-Enhancing Technique	Federated Learning
Privacy Guarantee	Data kept on premises





Dishonesty proportion

Secure Aggregation

Differential Privacy



Dishonesty proportion

Secure Aggregation

Differential Privacy



Secure Aggregation

Differential Privacy





Secure Aggregation

Differential Privacy



Secure Aggregation

Differential Privacy

Need for Lotto

Federated Learning

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Secure Aggregation

Differential Privacy

Need for Lotto

Population $(|0^4 - |0^8)$



Secure Aggregation

Differential Privacy

Need for Lotto

Population $(|0^4 - |0^8)$ **Selected participants** $(|0^2 - |0^3)$



Secure Aggregation

Differential Privacy

Need for Lotto



• **Random**: uniform chance



Secure Aggregation

Differential Privacy

Need for Lotto

Population $(|0^4 - |0^8)$ **Selected participants** $(|0^2 - |0^3)$

- **Random**: uniform chance
- **Informed**: "best-performing" clients are preferred (e.g., high speed and/or rich data)



Secure Aggregation

Differential Privacy

Need for Lotto



Problem: participant selection can be manipulated by the malicious server

Lotto - Overview

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No peer-to-peer network: all traffic relayed by the server

Threat model: malicious server colluding with some clients, and a public key infrastructure (**PKI**)

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Functionality

Support both **random** and informed selection

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Support both **random** and informed selection

Theoretical guarantee of

preventing manipulation

Lotto - Overview

Security

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Functionality

Support both **random** and informed selection

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Lotto - Overview

Security

Efficiency

Mild runtime overhead with no **network cost**

Problem: Random selection

Problem: Random selection



Selection criteria: <3

Problem: Random selection

Curr roun	ent d: 2	5
	Randomness	Select
#	$RF_{pkl}(2) = 9$	No
#2	$RF_{pk2}(2) = 1$	Yes
#3	$RF_{pk3}(2) = 7$	No

Selection criteria: <3




Selection criteria: <3 For dishonest majority

Problem: Random selection

Potential approach:

• Outcome verification





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1	3?	
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Selection criteria: <3 For dishonest majority

Problem: Random selection

Potential approach:

- Outcome verification
- Only within participants (10² 10³)









What is achieved:

Each participant sees a list of peers

Potential approach:

- Randomness verification
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What is achieved: Each participant sees a list of peers who presents only by chance.

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Output range: [0, 10)

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Problem: The server may arbitrarily **ignore honest** clients



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Unbounded advantage in growing dishonesty

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Solution: Enforce a large enough list and a small enough chance.



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Example

- **len(list)**: ≥ 200
- Chance: $\leq 0.1\%$



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► 0.99 .0 Pr. Fail in Half Dishonesty Example • **len(list)**: ≥ 200 0.5 • **Chance**: $\leq 0.1\%$ 0.0 80000 100000 120000 # Dishonest clients 76 Selected ≤ **50%** ≥ **50%**





What is achieved: Predictable to server? Each participant sees a list of peers who presents only by chance.



Public Round index **Examples**: #2 will be selected as $\mathbf{RF}_{pk2}(2) = 1 < 3$. Public Public keys

What is achieved: Predictable to server? Each participant sees a list of peers who presents only by chance.



Problem: Attack surfaces **enlarged**!

Examples: #2 will be selected as $\mathbf{RF}_{pk2}(2) = 1 < 3$. It's honest, so the server may grow its advantage by



Focused h	nacking
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Problem: Attack surfaces **enlarged**!

Examples: #2 will be selected as $\mathbf{RF}_{pk2}(2) = 1 < 3$.



What is achieved:PredictableEach participantto server?sees a list of peers who)presents only by chance.

The absent will not get arbitrarily ignored

¹Micali et al. "Verifiable random functions", In FOCS '99 ²Dodis et al. "A verifiable random function with short proofs and keys", In PKC '05 Solution: Self-sampling with

verifiable random functions (VRFs)^{1,2}.



Evaluation: **VRF.eval**_{sk2}(2) = (I,) (output, Secret key —



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Actual participants throughout the training?

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Problem: The server may **not follow**.

Involve non-selected dishonest ones



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Disregard **selected honest** ones



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arbitrarily ignored

¹Thus also of distributed DP (other privacy-enhancing techniques may not have this feature and this is left for future work).

Solution: Utilize existing secure semantics of secure aggregation¹

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• **Commitment**: necessary info shared only once



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Minor issues:

. . .

- Fixed sample size: over-selection
- Consistent round index: uniqueness check

Please find more in the paper :)

Problem: Informed selection

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Problem: Informed selection



Selection criteria: the fastest For dishonest majority



Problem: Informed selection

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Problem: Informed selection

Major Challenge: Client metrics are hard to verify by honest clients



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Major Challenge: Client metrics are hard to verify by honest clients

Metrics are fake





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Metrics are true, but...





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Major Challenge: Client metrics are hard to verify by honest clients

Metrics are fake



Metrics are true, but...



Solution: Approximate inform selection by **random** selection

Please find more in the paper :)

What can be **proven**:



Population





Population

Base rate of dishonest clients



What can be **proven**:



Base rate of dishonest clients



Example

- **Population**: 200,000
- Dishonesty base rate: 0.005

What can be **proven**:

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Base rate of dishonest clients



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- **Population**: 200,000
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- Target participants: 200

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Base rate of dishonest clients



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¹Random selection as an example. See results for informed selection in the paper.



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Oort^I → State-of-the-art **informed** selector: optimized for **time-to-accuracy** of training

Lai et al. "Oort: Efficient Federated Learning via Guided Participant Selection", In OSDI '21

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Lotto: Results summary



Support both random (exact) and informed (well **approximated)** selection



Theoretical guarantee (tight probability bound) of preventing manipulation



github.com/SamuelGong/Lotto

Security

Mild runtime overhead (≤10%) with no network cost (<1%)

Efficiency

Thank you

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